Games of No Strategy and Low-Grade Combinatorics

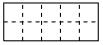
James Propp (jamespropp.org), UMass Lowell

Mathematical Enchantments (mathenchant.org)

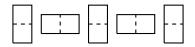
presented at MOVES 2015 on August 3, 2015

Slides at http://jamespropp.org/moves15.pdf

Boring puzzle: Tile a 2-by-*n* rectangle (here n = 5)



with n dominos.



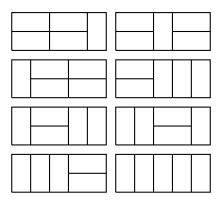
Interesting problem: In how many ways can a 2-by-n rectangle be tiled by dominos?

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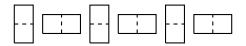
Answer (old result; whose?): The *n*th Fibonacci number.



Boring puzzle: Tile an "Aztec diamond of order n" (here n = 2)



with n(n+1) dominos.



(Aside: An Aztec diamond of order *n* consists of rows of length 2, 4, 6, ..., 2n - 2, 2n, 2n - 2, ..., 6, 4, 2, all centered.)

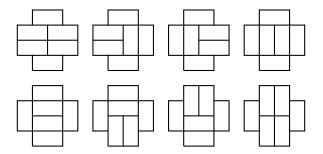
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Grades of enumerative combinatorics

These results are part of the branch of mathematics called **enumerative combinatorics**.

I like to split the subject into "grades", according to the growth-rates of the functions involved.

I say f(n) is of **grade** k if f(n) grows like exp n^k up to smaller correction factors; more precisely, if $(\log \log f(n)) / \log n \to k$.

The sequence of Fibonacci numbers is of grade 1; the sequence whose *n*th term is $2^{n(n+1)/2}$ is of grade 2.

0,1,2,...?

Combinatorics of grade 0 (e.g., figurate numbers) is ancient.

Combinatorics of grade 1 (e.g., Fibonacci numbers, Catalan numbers) goes back several centuries.

Combinatorics of grade 2 is just over a century old.

Aside from a theorem of Linial's on generalized spanning trees in hypergraphs, and a handful of related results, enumerative combinatorics has not progressed beyond grade 2.

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Is there such a thing as enumerative combinatorics of grade 3?

If so, how can we "graduate" to it?

Let's return to these questions later.

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Legal move: Join two free ends with a curve that does not cross any previously drawn curve (including the original circle) and then put a short stroke across the curve to create two new free ends, one on either side of the new curve. Delete the old free ends.

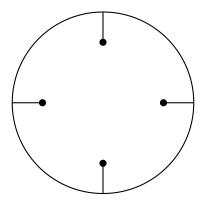
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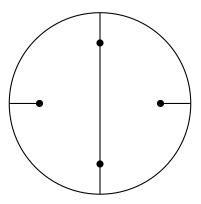
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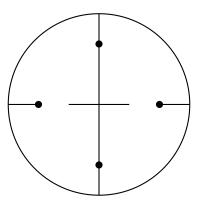
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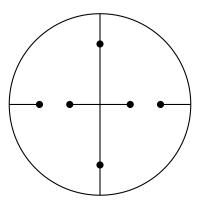
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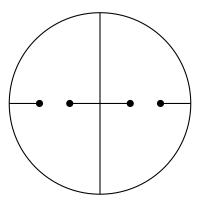
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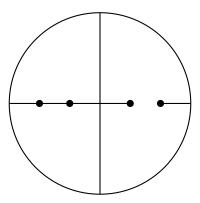




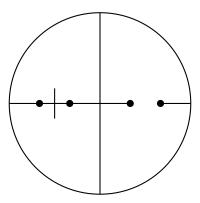




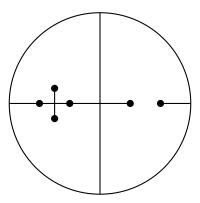
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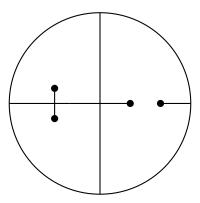
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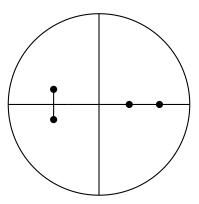
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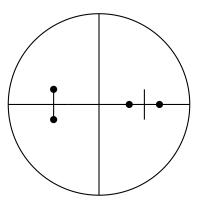


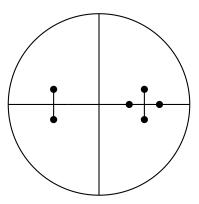
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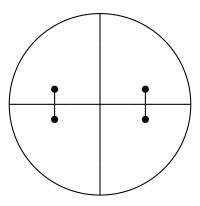
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Example (n = 4): First player moves:



First player wins!

Interesting problem: How short or long can a game be, and how many lines of play are there?

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Question: Is there an exact formula for the number of end-positions of the game?

Boring game: A chip firing game

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Initial state: A row of 2n + 1 bins. The middle bin has 2 chips. Every other bin contains 1 chip, except for the two outermost bins, which are empty.

 $0 \ 1 \ 1 \ \dots \ 1 \ 1 \ 2 \ 1 \ 1 \ \dots \ 1 \ 1 \ 0$

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Note that this number is highly composite: it is on the order of exp n^2 , but all its prime factors are less than n^2 .

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Enumerative combinatorics guild secret: These games are just square standard Young tableaux in disguise.

There's a general theory of chip-firing games, and it predicts that under very general conditions, the number of moves is predestined. But nobody has looked at the number of lines of play.

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The resulting sequence

 $1, 1, 1, 2, 4, 252, 2304, 343712160, 17361257184, \ldots$

just got added to the OEIS a few days ago.

"Welcome to the third grade!" (?)

What makes this sequence so intriguing is a combination of properties:

(1) It's of the THIRD grade.

(2) The terms have lots of small prime factors I can't explain.

 $\begin{array}{l} 252 = 2^2 \times 3^2 \times 7^1 \\ 2304 = 2^8 \times 3^2 \\ 34371260 = 2^5 \times 3^4 \times 5^1 \times 11^1 \times 2411^1 \\ 17361257184 = 2^5 \times 3^2 \text{ times a big prime} \\ \text{the next term} = 2^7 \times 3^1 \times 11^1 \times 13^1 \times 79^1 \text{ times a big prime} \\ \text{the next term} = 2^4 \times 5^1 \times 17^1 \times 43^1 \times 97^1 \text{ times a product of} \\ \text{three big primes} \end{array}$

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Does the product of all (or some) of the small primes have some combinatorial meaning?

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Statistical physicists have found exact solutions to lots of lattice models (such as tiling models) in 1D and 2D, but not 3D.

Perhaps we could break this dimension barrier if we knew more about grade 3 combinatorics.

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Winner: Last player to make a legal move

That is, the player who uses up all remaining score-lines wins.

Example: From

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L	1	1	
1	1	1	

a player can move to

$$\begin{array}{c} \cdot & \cdot \\ \cdot & \cdot \\ \cdot & \cdot \\ \cdot & \cdot \\ \cdot & \cdot \end{array} \text{ or } \begin{array}{c} - & - \\ - & - \\ \cdot & - \\ \cdot & \cdot \end{array} \text{ or } \begin{array}{c} - \\ - & - \\ \cdot & - \\ \cdot & - \end{array} \text{ or } \begin{array}{c} - \\ - & - \\ \cdot & - \\ \cdot & - \end{array}$$

In the first case, we use up 3 lengths of score-line; in the second and third cases, we use up 2 lengths of score-line.

Ignore this slide!

Claim: If a or b is even, a-by-b is a win for the first player.

Proof: Go first, and divide the rectangle into two identical pieces; thereafter, mimic your opponent's move.

You can abuse yourself by working out a winning strategy for the second player when a and b are both odd.

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Proof: Go first, and divide the rectangle into two identical pieces; thereafter, mimic your opponent's move.

You can abuse yourself by working out a winning strategy for the second player when a and b are both odd.

But you shouldn't, because it turns out that EVERY move is a win!

To see why, it's best to ignore everything I've told you in the last two minutes, because it was all designed to distract you from what's really going on.

But don't ignore this one

There's a one sentence-proof that the duration of the game is independent of the moves that are made.

You might want to work on this during the break. If you can't solve the problem, ask a child.

As for the number of lines of play: High school students Caleb Ji, Robin Park, and Angela Song, under the supervision of Tanya Khovanova, and with assistance from Pavel Etingof, have studied the case of a 2-by-n bar.

Letting B_n denote the number of lines of play, they showed that if p is 2, 5, or a prime that is congruent to 1 or 4 mod 5, then

 B_n is divisible by p for all sufficiently large n, whereas if p is ANY other prime, then this is NOT the case.

Further reading

R.J. Anderson, L. Lovász, P.W. Shor, J. Spencer, E. Tardos, and S. Winograd, Disks, balls, and walls: analysis of a combinatorial game, American Mathematical Monthly, Volume 96, Number 6, June-July 1989, pages 481–493.

Grant Cairns and Korrakot Chartarrayawadee, Brussels Sprouts and Cloves, Mathematics Magazine, Volume 80, Number 1, February 2007, pages 46–58(13).

Richard Guy, She Loves Me, She Loves Me Not: Relatives of two games of Lenstra, https://oeis.org/A006016/a006016_1.pdf.

Gil Kalai, Enumeration of Q-acyclic simplicial complexes, Israel J. Math. 45 (1983), no. 4, 337351.

James Propp, Games of no strategy, in preparation.

Slides at http://jamespropp.org/moves15.pdf \Box , $d \in \mathbb{R}$, $d \in \mathbb{R}$, $d \in \mathbb{R}$